

## 2010 AwesomeMath UTD Team Contest Part 2

### 1 Algebra

1. Let  $n \geq 2$ . How many polynomials  $Q(x)$  with real coefficients of degree at most  $n - 1$  are there such that

$$x(x-1)\cdots(x-n)Q(x) + x^2 + 1$$

is the square of a polynomial?

**Solution** The given condition says

$$f(x)^2 = x(x-1)\cdots(x-n)Q(x) + x^2 + 1 \quad (1)$$

for some polynomial  $f(x)$  of degree at most  $n$ . Plugging  $x = 0, 1, \dots, n$  into (1) gives

$$f(x) = \pm\sqrt{x^2 + 1}, \text{ when } x = 0, 1, \dots, n. \quad (2)$$

The following is key:

*Fact 1.* Given  $n + 1$  points  $(x_0, y_0), \dots, (x_n, y_n)$  with distinct  $x$ -coordinates, there exists exactly one polynomial  $f$  of degree at most  $n$  so that  $f(x_i) = y_i$  for  $i = 0, 1, \dots, n$ .

Applying this to (2) we get  $2^{n+1}$  possibilities for  $f(x)$  since we have 2 choices of sign for each of  $x = 0, 1, \dots, n$ . If  $f(x)$  is a solution to (2) then so is  $-f(x)$ ; we get  $2^n$  possibilities for  $f(x)^2$ . Solve (1) to get  $2^n$  possibilities for  $Q(x)$ :

$$Q(x) = \frac{f(x)^2 - x^2 - 1}{x(x-1)\cdots(x-n)}$$

Each such polynomial is a valid solution because  $f(x)^2 - x^2 - 1$  is zero at  $x = 0, 1, \dots, n$  and hence is divisible by  $x(x-1)\cdots(x-n)$ . ■

2. Let  $a_1, \dots, a_5$  be real numbers and  $x, y$  be real numbers such that

$$(a_1 + a_2 + a_3 + a_4 - a_5)^2 \geq 3(a_1^2 + a_2^2 + a_3^2 + a_4^2 - a_5^2).$$

Prove that

$$(a_1 + a_2 + a_3 + a_4 - a_5 - x - y)^2 \geq a_1^2 + a_2^2 + a_3^2 + a_4^2 - a_5^2 - x^2 - y^2.$$

**Solution** By  $T_2$ 's Lemma (or Cauchy-Schwarz, or QM-AM),

$$\begin{aligned} \frac{(a_1 + a_2 + a_3 + a_4 - a_5 - x - y)^2}{1} + \frac{x^2}{1} + \frac{y^2}{1} &\geq \frac{(a_1 + a_2 + a_3 + a_4 - a_5)^2}{3} \\ &\geq a_1^2 + a_2^2 + a_3^2 + a_4^2 - a_5^2. \end{aligned}$$

This rearranges to the desired inequality.

3.  $n > 1$  bunnies sit on a number line such that the maximum distance between any two of them is  $d$ . Each step, 2 bunnies are selected. The bunny at the left, say  $A$ , jumps some distance  $x > 0$  to the right, while the bunny at the right, say  $B$ , jumps the same distance  $x$  to the left, such that  $A$  is still to the left of (or occupies the same location as)  $B$ . (Bunnies may jump over each other.) After a finite number of steps, let  $S$  be the sum of the distances traveled by all bunnies. Let  $L(n, d)$  be the smallest number so that we always have  $S \leq L(n, d)$ . Find  $L(n, d)$ .

**Solution** Number the bunnies  $1, 2, \dots, n$ , and let their positions be  $x_1, \dots, x_n$ , respectively, where  $0 \leq x_1, \dots, x_n \leq d$ . Define

$$f(x_1, \dots, x_n) = \sum_{1 \leq i < j \leq n} |x_i - x_j|. \quad (3)$$

*Lemma 1.* After the bunnies travel a distance of  $l$ , then  $f(x_1, \dots, x_n)$  decreases by at least  $l$ . Equality holds iff the only time bunnies move jump toward each other is when there are no other bunnies in between.

*Proof.* Decomposing each step into substeps as necessary, we may assume that no bunny crosses another bunny during a step. It suffices to prove the claim for a step. WLOG, assume the bunnies are in increasing order from left to right. Suppose bunnies  $i$  and  $j$  at positions  $x_i < x_j$  each move a distance of  $l$  during this step. Terms in (3) not involving  $x_i, x_j$  do not change. Now, for  $k < i$  or  $k > j$ , we can pair the terms  $|x_k - x_i| + |x_k - x_j|$ ; one increases by  $l$  while the other term decreases by  $l$ . If  $i < k < j$ , then both terms  $|x_k - x_i| + |x_k - x_j|$  decrease by  $l$ . The term  $|x_j - x_i|$  decreases by  $2l$ , so the whole sum decreases by at least  $2l$ , with equality iff there is no bunny between bunnies  $i$  and  $j$ .  $\square$

Hence no the total distance the bunnies move cannot be greater than  $f(x_1, \dots, x_n)$ , where  $x_1, \dots, x_n$  are the initial positions. The function  $f(x_1, \dots, x_n)$  achieves a maximum on the closed and bounded set  $0 \leq x_1, \dots, x_n \leq d$ ; since the functions  $|x - a|$  are all convex, by looking at  $f$  in each variable separately, we get that the maximum must be obtained when each variable  $x_1, \dots, x_n$  is equal to 0 or  $d$ . If  $k$  of them are equal to 0, then

$$f(x_1, \dots, x_n) = k(n - k)d,$$

which attains maximum for the value of  $k$  closest to the vertex  $\frac{n}{2}$  of the quadratic  $x(n - x)$ . When  $n$  is even the maximum is  $\frac{n^2 d}{4}$ ; when  $n$  is odd the maximum is  $\frac{(n-1)(n+1)d}{4}$ . Now we show these are indeed the values of  $L(n, d)$ .

*Lemma 2.* Given any arrangement of bunnies in increasing order from left to right, if we repeatedly move pairs of bunnies  $(1, 2), (2, 3), \dots, (n-1, n)$ , each time moving the 2 bunnies in the pair to the same location, in that order, then the location of all bunnies will converge to  $\bar{x} = \frac{\sum_{i=1}^n x_i}{n}$ .

*Proof.* Since the sum of the positions of the bunnies is invariant, if the position of all bunnies converges to  $x$  then we have  $nx = x_1 + \dots + x_n \Rightarrow x = \bar{x}$ . Let  $x_i(t)$  be the position of the  $i$ th bunny after  $t$  repetitions of the algorithm above. Then  $x_n(t) - x_1(t)$  is decreasing, so it converges to some  $D \geq 0$ . It suffices to consider  $D \neq 0$ . Note that after the  $(t+1)$ th step of the algorithm, every point between  $x_1(t)$  and  $x_n(t)$  has been jumped over (or on) by one bunny. Indeed, after moving pairs  $(1, 2), \dots, (i-1, i)$  to the same location, bunny  $i$  has not moved right. Then when  $(i, i+1)$  are moved together, all points between  $x_i(t)$  and  $x_{i+1}(t)$  have been jumped over (or on). So the total distance traveled by the bunnies in the  $(t+1)$ th step is at least  $x_n(t) - x_1(t) \geq D$ , and the total distance traveled by the bunnies is unbounded if the step is repeated, a contradiction.  $\square$

Now if  $n$  is even, then take  $x_1 = \dots = x_{\frac{n}{2}} = 0$  and  $x_{\frac{n}{2}+1} = \dots = x_n = d$ , and if  $n$  is odd, take  $x_1 = \dots = x_{\frac{n+1}{2}} = 0$ ,  $x_{\frac{n+3}{2}} = \dots = x_n = d$ , and carry out the algorithm described in the lemma. Then by Lemma (1), in the limit all the bunnies travel a total distance of  $\frac{n^2d}{4}$  or  $\frac{(n^2-1)d}{4}$ , respectively, because we have that  $f$  attains this value, and no bunny jumps over another bunny. Hence

$$L(n, d) = \begin{cases} \frac{n^2d}{4}, & n \text{ even} \\ \frac{(n^2-1)d}{4}, & n \text{ odd} \end{cases} \quad (4)$$

■

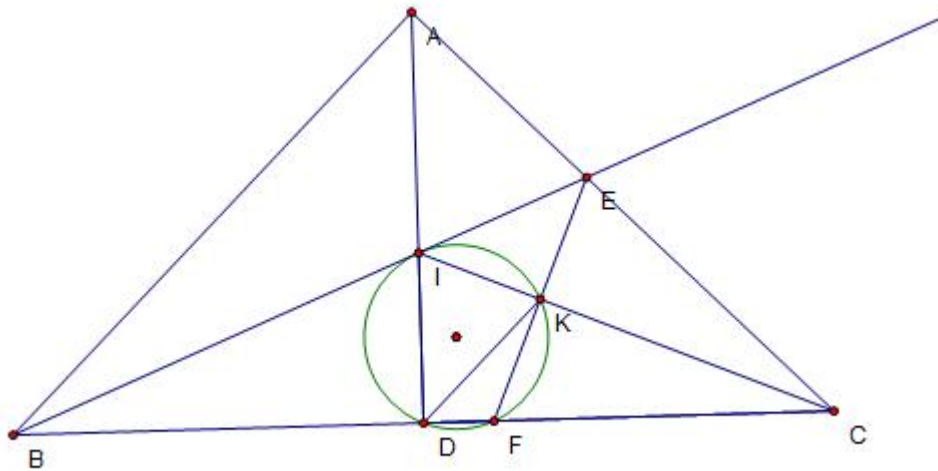
## 2 Combinatorics

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## 3 Geometry

1. Let  $ABC$  be a triangle with  $AB = AC$ . The angle bisectors of  $\angle CAB$  and  $\angle ABC$  meet the sides  $BC$  and  $CA$  at  $D$  and  $E$ , respectively. Let  $K$  be the incenter of triangle  $ADC$ . Suppose that  $\angle BEK = 45^\circ$ . Find all possible values of  $\angle CAB$ .

**Solution**  $60^\circ$  or  $90^\circ$ .



The angle bisectors  $AD$  and  $BE$  intersect at the incenter  $I$  of  $\triangle ABC$ . Reflect  $E$  across  $CI$  to point  $F$  on  $BC$ . Then  $EF \perp CI$  and  $\angle IFK = \angle IEK = 45^\circ$ . Since  $K$  is the incenter of  $\triangle ADC$ ,  $DK$  bisects  $\angle ADC$ , and  $\angle IDK = 45^\circ = \angle IFK$ . Hence  $I, K, D, F$  are concyclic.

If  $F = D$  then  $CE = CF = CD$ . Since  $AD \perp BC$  and  $BE, AC$  are the reflections of  $AD, BC$  across  $CI$ , we get  $BE \perp AC$  and  $AB = AC$ , i.e.  $\triangle ABC$  is equilateral and  $\angle CAB = 60^\circ$ .

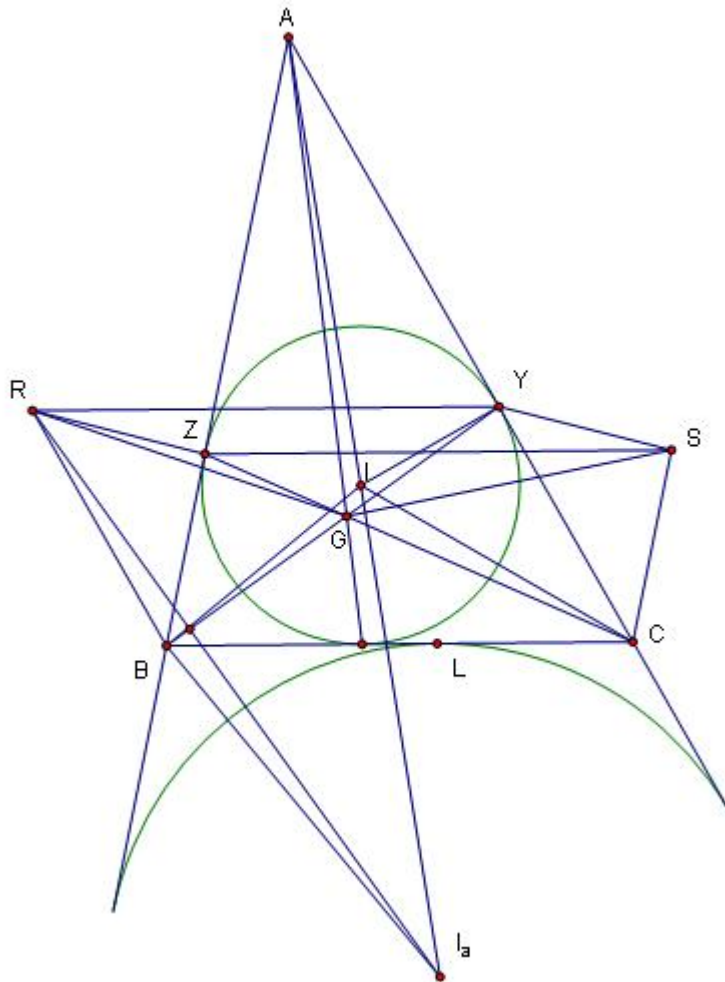
Else,  $\angle IKE = \angle IKF = \angle IDF = 90^\circ$ , and

$$\begin{aligned} \angle A &= \angle BEC - \angle EBA = 45^\circ + \angle KEC - \angle EBA \\ &= 45^\circ + \angle EKI - \angle ECI - \angle EBA = 45^\circ + 90^\circ - (90 - \angle A/2) \end{aligned}$$

and  $\angle A = 90^\circ$ . ■

- Let  $ABC$  be a triangle. The incircle of  $ABC$  touches the sides  $AB$  and  $AC$  at the points  $Z$  and  $Y$ , respectively. Let  $G$  be the point where the lines  $BY$  and  $CZ$  meet, and let  $R$  and  $S$  be points such that the two quadrilaterals  $BCYR$  and  $BCSZ$  are parallelograms. Prove that  $GR = GS$ .

**Solution**



Let the excircle  $\omega_A$  opposite  $A$  intersect  $AB$  and  $AC$  at  $D$  and  $E$ , respectively, let  $a, b, c$  be the side lengths of the triangle and let  $s$  be the semiperimeter. Now,

$$BR = CY = s - c = BD,$$

so  $B$  has equal power with respect to  $\omega_A$  and the point (degenerate circle)  $R$ . We have

$$YR = CB = (s - b) + (s - c) = CY + CE = YE$$

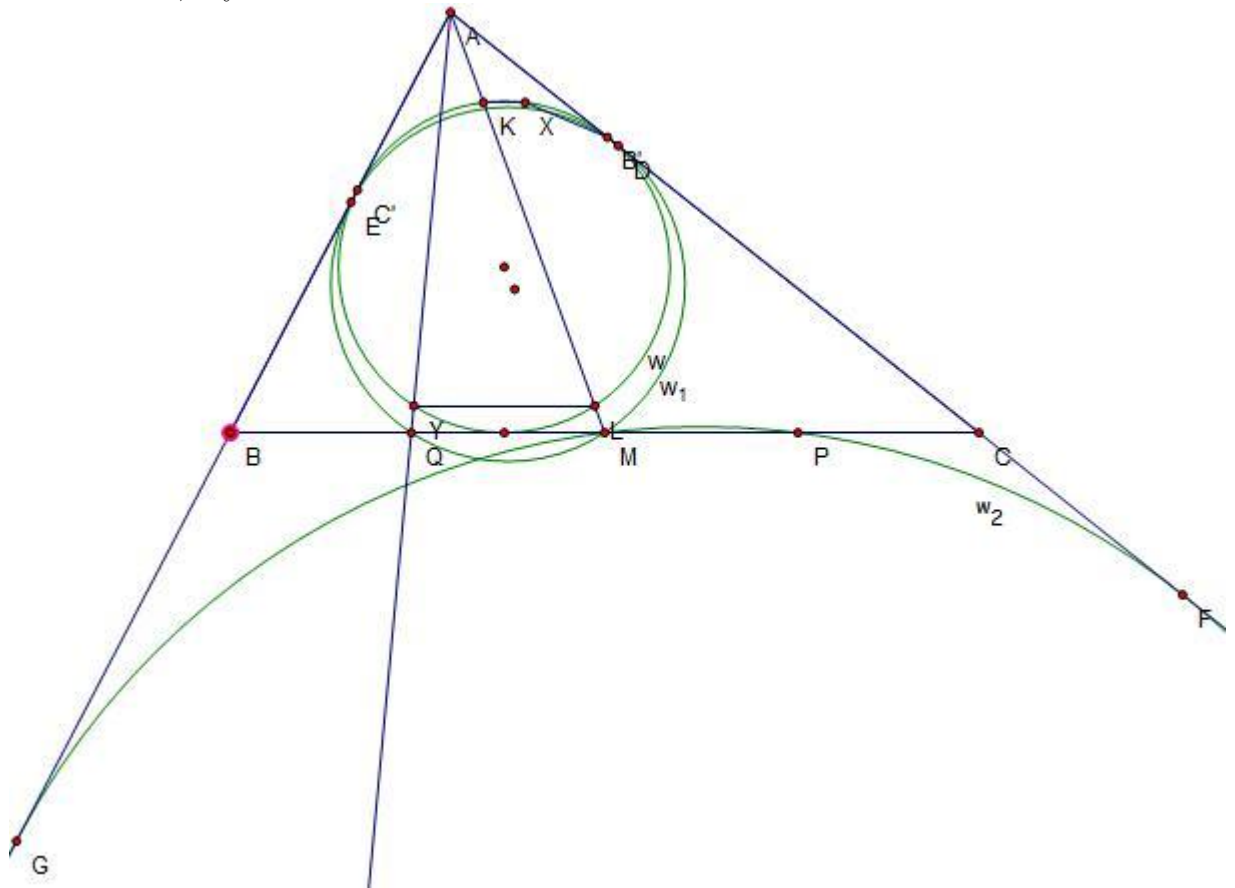
so  $Y$  has equal power with respect to  $\omega_A$  and  $R$ . Hence  $BY$  is the radical axis of  $R$  and  $\omega_A$ . Similarly,  $CZ$  is the radical axis of  $S$  and  $\omega_A$ . Hence their point of intersection  $G$  is the radical center of  $R, Y$ , and  $\omega_A$ . Hence  $G$  has equal power with respect to  $R$  and  $S$ , and  $GR = GS$ . ■

- The median  $AM$  of  $\triangle ABC$  intersects its incircle  $\omega$  at  $K$  and  $L$ . The lines through  $K$  and  $L$  parallel to  $BC$  intersect  $\omega$  again at  $X$  and  $Y$ . The lines  $AX$  and  $AY$  intersect  $BC$  at  $P$  and  $Q$ . Prove that  $BP = CQ$ .

**Solution** Without loss of generality,  $AC > AB$ .

Let  $\omega$  intersect  $AB$  and  $AC$  at  $C'$  and  $B'$ , respectively. Since  $YL \parallel QM$ , there is a homothety  $H_1$  centered at  $A$  sending  $YL$  to  $QM$ .  $H_1$  sends  $\omega$  to a circle  $\omega_1$  tangent to  $AB$  and  $AC$ . Suppose  $\omega_1$  intersects  $AB$  and  $AC$  at  $E$  and  $D$ . Similarly,

a homothety  $H_2$  centered at  $A$  sends  $KX$  to  $MP$ , and  $\omega$  to a circle  $\omega_2$  tangent to  $AB$  and  $AC$ , say at  $G$  and  $F$ .



By equal tangents,

$$EG = AG - AE = AF - AD = DF. \tag{5}$$

By power-of-a-point with respect of  $B$  and  $C$  with respect to  $\omega_1$  and  $\omega_2$ , setting  $BQ = x, CP = y, BC = a$  we get

$$\begin{aligned} BE^2 &= BQ \cdot BM = x \cdot \frac{a}{2} \\ BG^2 &= BM \cdot BP = \frac{a}{2} \cdot (a - y) \\ CD^2 &= CM \cdot CQ = \frac{a}{2} \cdot (a - x) \\ CF^2 &= CP \cdot CM = y \cdot \frac{a}{2} \end{aligned}$$

From (5), we get  $BE + BG = CD + CF$ . Substituting the above into this equation and dividing by  $\sqrt{a/2}$ ,

$$\begin{aligned} \sqrt{x} + \sqrt{a - y} &= \sqrt{a - x} + \sqrt{y} \\ \sqrt{x} - \sqrt{a - x} &= \sqrt{y} - \sqrt{a - y} \end{aligned}$$

Since  $\sqrt{x} - \sqrt{a - x}$  is an increasing function,  $x = y, BQ = CP, BP = CQ$ . ■

## 4 Number Theory

1. There are  $n \geq 51$  points in the plane with integer coordinates, such that the distance between any two is an integer. Prove that at least 49 percent of the distances are even.

**Solution** Let the points be  $(x_1, y_1), \dots, (x_n, y_n)$ . We show that either all the  $x_i$  have the same parity, or all the  $y_i$  have the same parity.

Suppose that the  $x_i$  do not all have the same parity; suppose  $x_i \not\equiv x_j \pmod{2}$ . Then the square of the distance between  $(x_i, y_i)$  and  $(x_j, y_j)$  is  $(x_i - x_j)^2 + (y_i - y_j)^2$ . If  $y_i \not\equiv y_j \pmod{2}$ , then  $(x_i - x_j)^2 + (y_i - y_j)^2 \equiv 1 + 1 \equiv 2 \pmod{4}$  so cannot be a square. Hence  $y_i \equiv y_j \pmod{2}$ . If  $x_i \equiv x_k \pmod{2}$  then finding  $x_j \not\equiv x_i \pmod{2}$  we get  $y_i \equiv y_j \equiv y_k \pmod{2}$  as well. Thus all the  $y_i$  have the same parity.

Without loss of generality, all the  $y_i$  have the same parity. Suppose  $k$  of the  $x_i$  them are even; then  $n - k$  of them are odd.  $(x_i, y_i)$  and  $(x_j, y_j)$  have even distance apart if  $x_i - x_j$  and  $y_i - y_j$  are both even. There are

$$\frac{k(k-1)}{2} + \frac{(n-k)(n-k-1)}{2} = \frac{k^2 + (n-k)^2 - n}{2}$$

such pairs. The proportion of even distances is at least

$$\frac{[k^2 + (n-k)^2 - n]/2}{n(n-1)/2} \geq \frac{\binom{n}{2} - n}{n(n-1)} = \frac{\frac{n}{2} - 1}{n-1} \geq \frac{\frac{51}{2} - 1}{51-1} = \frac{49}{100}.$$

■

2. Let  $p, q, r$  be distinct primes such that

$$pq \mid r^p + r^q.$$

Prove that either  $p$  or  $q$  equals 2.

**Solution** Suppose the relation holds but  $p \neq 2, q \neq 2$ . By Fermat's Little Theorem,  $r^p \equiv r \pmod{p}$  and  $r^q \equiv r \pmod{q}$ . Then since  $r$  is relatively prime to  $p, q$ ,

$$\begin{aligned} r^p + r^q &\equiv 0 \pmod{p} \implies \\ r^{q-1} &\equiv -1 \pmod{p} \\ r^p + r^q &\equiv 0 \pmod{q} \implies \\ r^{p-1} &\equiv -1 \pmod{q} \end{aligned}$$

Since  $-1 \not\equiv 1 \pmod{p, q}$ , we get

$$\text{ord}_p(r) \nmid q-1, \text{ord}_q(r) \nmid p-1. \tag{6}$$

Since

$$\begin{aligned} r^{2(q-1)} &\equiv 1 \pmod{p} \\ r^{2(p-1)} &\equiv 1 \pmod{q}, \end{aligned}$$

we get

$$\text{ord}_p(r) \mid 2(q-1), \text{ord}_q(r) \mid 2(p-1). \quad (7)$$

For an integer  $n$  let  $v_2(n)$  denote the highest power of 2 dividing  $n$ . Let  $x = v_2(\text{ord}_p(r))$  and  $y = v_2(\text{ord}_q(r))$ . From relations in (6) and (7),

$$\begin{aligned} x &= v_2(2(q-1)) = v_2(q-1) + 1 \\ y &= v_2(p-1) + 1. \end{aligned} \quad (8)$$

By Fermat's Little Theorem,  $\text{ord}_p(r) \mid p-1$  and  $\text{ord}_q(r) \mid q-1$ . Hence

$$\begin{aligned} x &\leq v_2(p-1) \\ y &\leq v_2(q-1) \end{aligned} \quad (9)$$

Putting (8) and (9) together, we get  $x \leq y-1, y \leq x-1$ , contradiction.  $\blacksquare$

3. Find all solutions in positive integers to  $a^2 + 2b^2 = (x^5 - x^3 - 1)(x^5 - x^3 - 3)$ . (Hint: You may use the fact that  $-2$  is a quadratic residue modulo a prime  $p$  if and only if  $p \equiv 2, 1, \text{ or } 3 \pmod{8}$ .)

### Solution

*Lemma 3.* Let  $a, b$  be integers. If  $p$  is an  $8k+5$  or  $8k+7$  prime, then the highest power of  $p$  dividing  $a^2 + 2b^2$  is an even power.

*Proof.* Suppose  $a^2 + 2b^2 \equiv 0 \pmod{p}$ . Let  $p^r$  be the highest power of  $p$  dividing  $a^2 + 2b^2$ . Let  $a' = \frac{a}{p^r}$  and  $b' = \frac{b}{p^r}$ ; then  $a', b'$  are not both divisible by  $p$ . We have that  $a'^2 + 2b'^2 = \frac{a^2 + 2b^2}{p^{2r}}$ ; it suffices to show that  $p \nmid a'^2 + 2b'^2$ . Suppose by way of contradiction that  $a'^2 + 2b'^2 \equiv 0 \pmod{p}$ . Since  $a'$  and  $b'$  are not both divisible by  $p$ , neither is. We get  $(a'b'^{-1})^2 \equiv -2 \pmod{p}$ , contradicting the fact that  $-2$  is not a quadratic residue modulo  $p$ .  $\square$

Note  $a = b = x = 1$  is a solution. We show this is the only solution. If  $x > 1$ , then both  $x^5 - x^3 - 1, x^5 - x^3 - 3 > 0$ . Now  $x^5 - x^3 \equiv 0 \pmod{8}$  so

$$\begin{aligned} x^5 - x^3 - 1 &\equiv 7 \pmod{8} \\ x^5 - x^3 - 3 &\equiv 5 \pmod{8} \end{aligned}$$

However, the product of an even number of  $8k+5, 8k+7$  primes and an arbitrary number of  $8k+1, 8k+3$  primes cannot be congruent to 5 or 7 modulo 8. Moreover,  $x^5 - x^3 - 1$  and  $x^5 - x^3 - 3$  are relatively prime since their difference is 2 and both are odd. Hence they have no common prime factor. It follows that  $(x^5 - x^3 - 1)(x^5 - x^3 - 3)$  has a  $8k+5$  or  $8k+7$  prime appearing an odd number of times, and cannot equal  $a^2 + 2b^2$ .  $\blacksquare$



## 5 Problem Credits

A1 Holden Lee

A2 Romania 1983

A3 Holden Lee

C1 Temporarily withheld.

C2 Temporarily withheld.

C3 Temporarily withheld.

G1 IMO 2009/4

G2 IMO Shortlist, 2009/G3 (Figure due to livetolove212)

G3 IMO Shortlist, 2005/G6

N1 Andrei Ciupan

N2 Based off many past Olympiad problems of the same flavor.

N3 Holden Lee

And... Holden Lee for L<sup>A</sup>T<sub>E</sub>Xing up everything!